Mobility Management: Deployment and Adaptability Aspects Through Mobile Data Traffic Analysis

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Abstract

The expected boost in mobile data traffic and the evolution towards the next generation of networks are making cellular operators reconsider whether current approaches for handling mobility could be improved, according to the characteristics of the mobile traffic that actually flows through real networks. In this work, we make use of extensive analysis of real network traces to infer the main characteristics of mobile data traffic for a particular operator. Our analysis focuses on the features related to mobility, i.e., location information, number of handovers, or duration of the data traffic exchange. New techniques to gather the mobility characteristics of the user based on data and control packets correlation are designed and applied to compare the gains of deploying different mobility and the network characteristics brings some benefits to the network operator over the current approach, especially in scenarios of low mobility, where a Distributed Mobility Management solution proves to be more efficient.

Keywords: Mobility, handover, Distributed Mobility Management (DMM), Big Data, Traffic Analysis

1. Introduction and Motivation

Cellular networks have evolved into extremely complex systems, where performance and behavior depend on the interaction of a multitude of logical modules. Coupled with this complexity, any analysis performed over these systems faces scalability challenges due to the size (in number of nodes) and amount of traffic served. The expected evolution of cellular networks forecasts an increase in both magnitudes, due to the growing traffic demands and the so-called radio access network (RAN) densification. In this scenario, operators struggle to monitor and analyze their networks through an amalgam of vendor specific probes and management systems, providing information which is difficult to aggregate and analyze. The lack of tools for the design and optimization of next generation networks, carrying several orders of magnitude more traffic and serving a wider set of possible clients (including machine and humans) is a challenge, requiring novel techniques that are able to provide trends, relations and design guidelines for the deployment of new systems. A promising trend in this area is the use of Big Data techniques to gather information on the behavior of the network, analyzing and

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inferring knowledge out of the myriad of data flows transported by the network [1]. It is a common practice to monitor the traffic flowing through the network to evaluate network performance or to look for delay or usage patterns [2]. With the increase of data traffic and the raise of more powerful processing techniques and capabilities, this traffic monitoring has turned to Big Data techniques to analyze the traffic [3]. In this work, we built on top of these tools to provide some insights on the usage of resources for mobility management with actual data from a real operator's network.

The work reported in this paper started as a discussion on how to analyze the benefits and drawbacks of new techniques for traffic offloading in a real operator environment, specifically if the application of mobility protocols based on the DMM concept were worth the effort considering the structure of the network, deployment characteristics and the real user traffic. In order to answer this question, we tried to directly analyze the data traffic in one of the operator's core interfaces looking for mobility patterns and traffic characteristics, required to understand the performance of the mobility protocol. Through this direct approach we found two main problems: i) The overwhelming number of flows going through the interface and their lack of information matching location of the user and ii) the need for a mechanism to correlate the information carried in the data path and the control path, to match flows to users and cells. To overcome these challenges, we resorted to the deployment and analysis of the data using big data

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analytic techniques. Therefore, this paper focuses i) on the development of a platform to perform Big Data analysis over packet traces in a real operator environment in Turkey under privacy and regulation concerns, and ii) to showcase the utility of this new framework by the analysis of the suitability of applying a new mobility management concept, DMM, to the current network deployment. In order to gain an understanding of the whole network operation, this work also reports on the challenges of working with packet captures taken on the data path of the operator's core, which can carry thousands of flows, and the need of correlating these flows with the control data obtained from a different interface in the operator's core.

Following these ideas, the remaining of the paper is structured as follows: Section 2 presents recent works in the area of Distributed Mobility Management and Traffic Analysis that are relevant for this work. Section 3 describes the method for the collection of the mobile data traffic information, which is analyzed in Section 3.4. Based on the results from these experimental measurements, in Section 4 we apply a DMM mechanism and compare it to the current deployment showcasing that it can be more convenient under certain mobility scenarios. Finally, Section 5 presents the final conclusions.

2. Related Work

Managing the mobile traffic is becoming increasingly complex, both technically and economically for mobile operators. This is not only because of increased data usage and emergence of new services and models, but also due to the change in users behaviours towards increased mobility. Mobility management support in cellular networks has always been a critical capability. Therefore, there has been numerous works outlining different mobility management approaches in the literature [4] [5] [6] [7] [8] [9]. A recent comprehensive tutorial on mobility management in data networks that especially focuses on network mobility and session migration is given in [4].

Current architectures for mobile and cellular networks are centralized [6] [7] and hierarchical[8] which drives the user traffic to go all the way up into the core network. Proxy Mobile IPv6 (PMIPv6) [7] and GPRS Tunneling Protocol (GTP) [9] protocols have been adapted for the applicability of the centralized mobility management solutions for mobile operators. Even though centralized mobility solutions reduce the signalling between the hosts and is relatively simple for following users movements, there are various issues with their adoption in cellular networks as identified in [10]. Some of these problems are: low scalability due to need for new mobility anchors as the number of mobile nodes and data traffic increases, per node mobility support that increases congestion on mobility anchors due to mobility support for all flows, single point of failure due to many mobile nodes connection to single mobility anchors [11], and non-optimal routes which may result in longer delays and excessive loads in core network.

As opposed to centralized solutions, distributed mobility management approach has proposed a flat and flexible architecture by ensuring per flow mobility support and optimal path routing in a highly scalable manner [4] [12] [10] [13]. A comprehensive overview of distributed mobility management techniques that includes standardization activities of both The 3rd Generation Partnership Project (3GPP) and Internet Engineering Task Force (IETF) has been given in [4]. The standardization activities for distributed mobility management are primarily driven by IETFs DMM working group [12]. The requirements for distributed mobility management has been recently given in Request for Comment (RFC) 7333 [10]. Some of these requirements constitute functionalities such as distributed processing, transparency to upper layers, IPv6 target deployments, re-usage of existing mobility protocols, co-existence with existing network deployments and hosts, by-passable mobility support for each applications and security considerations. From a practical implementation aspect, the authors in [13] have demonstrated the first practical evaluation results of distributed mobility management approach based on evaluations with real Linux-based prototype implementations.

Big data and data analytics are recently emerging to facilitate the development of new analytics applications, and to leverage the mobile operators understanding and exploitation of data which is constantly flowing through their networks infrastructures. In a more general point of view, one can find big data platform utilization (namely Hadoop [14]) and exploitation of data analytics by telecom operators in different recent studies [15] [16] [17] [18] [19]. For example, application areas of big data analysis applied by telecom operators ranges from anomaly detection for IT infrastructure security and resiliency [16], network coverage analysis [15], proactive caching for 5G [17] [18] to social network analysis for consumer behavior modeling [19].

However, none of the studies described above have demonstrated the deployment and adaptability aspects of distributed mobility management based solutions inside a real mobile operator by exploiting a comprehensive mobile data analysis. Therefore, it is clear that in order to embrace new technologies for long term alternatives to current centralized cellular infrastructures, recently proposed DMM solutions need to be further investigated well in the context of applicability and adaptability for mobile operators. Based on this observation, in this article, we focus on validation of DMM approach's performance results using a real operator's network data by exploiting big data techniques for data analysis. The data management (extraction, analysis and usage) process studied in this paper is specific to our distributed mobility management scenario and plays an important role in supporting the applicability of the proposed solution. As far as we know, the comparisons of distributed mobility management with current centralized mobility management approaches with large amount of mobile data usage is not available in the literature. Using the analysis in this paper, some interesting

conclusions are deduced by bringing real world considerations about the applicability of the new distributed mobility management solution inside mobile operator's network.

3. Description of the system

As explained in Section 1, this work aims at providing an analysis of the suitability of applying a new mobility management concept, DMM, to the current network deployment. In order to do so, we needed to develop a new system to be able to i) gather user data traces, ii) correlate the overwhelming set of user and control information to extract relevant mobility information and iv) use the mobility information extracted to apply an analytical model for the packet delivery cost of DMM and compare it with current deployments. In the following sections we explain each of the steps in this complex process.

3.1. Dataset and collection of information

Identifying the structural patterns in the data traffic is of high importance for mobile operators in order to apply optimal mobility management techniques within their network. Mobile network traffic has a highly complex and massive structure, making it tough to analyse and reveal structural patterns. It is not unusual that one or more terabytes of data per second is flowing in a typical mobile operator consisting of 10 to 20 million subscribers which translates into roughly exabytes every month. The scale of this problem rules out direct sniffing approaches [20], with the additional problem of data packets not carrying information regarding the location of the user. The analysis of mobility requirements necessitates the extraction of handover related information, which can be tracked both in radio access and core network nodes. Accessing handover related information from the radio access nodes is difficult since the amount of probes that needs to be placed into the infrastructure can be large. Moreover, log data from different entities, such as the radio network controller (RNC), can be too hard to extract for further mobility analysis due to the unavailability of appropriate tools, which are mostly vendor specific. In High Speed Packet Access (HSPA), the core network is only notified of Location Area (LA) and Routing Area (RA) updates, when the mobile terminal is in the idle state, since in connected state the access network can still locate it and report to the core network when necessary. Moreover, our approach aims to link user roaming across the network with the characteristics of the mobile traffic being exchanged.

Due to the aforementioned problems, we propose a system based on the extraction of handover related information observed in the core network nodes, where the correlation of the control message headers with the information of the flows in the user data plane are obtained by capturing the data and control packets in the Gn interface. One of the constraints imposed in the system design was the lack of existing measurement tools over the interfaces for

mapping flow information with the location information inside the operator domain. The current infrastructure allows capturing control packets in the 3GPP Gn interface between the Gateway GPRS Support Node (GGSN) and the Serving GPRS Support Node (SGSN) (see Fig. 1). Specifically, we track the Create Context and Update Context messages of the Packet Data Protocol (PDP) [9]. In this way, we can simplify the requirements for monitoring user mobility avoiding the need for multiple probes. For our analysis, we do not need to track the exact location of the user at the precise moment that it happens, but detect the changes in their point of attachment and characterize the traffic being exchanged by roaming users. We focus our analysis on handovers involving RA changes, which is inline with the current DMM architectures discussed in the IETF.¹

The final outcome of this process provides a combined and detailed listing of the control and user plane packets indicating location of mobile terminals via the information included in PDP *Create Context* and *Update Context* messages. This trace is further processed to characterise mobile data traffic (see Section 3.4) and to obtain insights on the applicability of the DMM concept in an operator's network, as explained in Section 4.

3.2. Architecture of the dataset collection mechanism

A general view of the architecture for the extraction of flow information is provided in Fig. 1. The system is composed of mainly two elements, the Flow Extraction Manager (FEM) and the Processing Cluster that has been implemented using Hadoop [14]. One of the existing Gn interfaces with high traffic in the core network between SGSN and GGSN is mirrored and collected in the FEM, which applies initial processing and transfers the data to be analysed and filtered into the Hadoop Cluster. The extracted flow information is sent back to FEM from the Hadoop Cluster for collecting analysis results. The analysis results are outputted by FEM for detailed observations of the network state. The details of the flow extraction process as well as the analysis in Hadoop cluster are given in Section 3.3. In the following, we detail some of the operations of the system.

3.2.1. GTP-U and GTP-C correlation on Gn Interface

Network packets sent from a user equipment (UE) to the packet data network (PDN), e.g. Internet, pass through the SGSN which tunnels them towards the GGSN. GTP is used for tunneling the packets in the Gn interface [9].

The operation of GTP protocol differentiates user and control planes. The user plane packets on the Gn interface flow on the GTP User (GTP-U) [9], which is in effect a relatively simple IP based tunneling protocol allowing several tunnels between each set of end points. When used in a

¹Note that routing areas may include one to several groups of cells connected to same RNC.

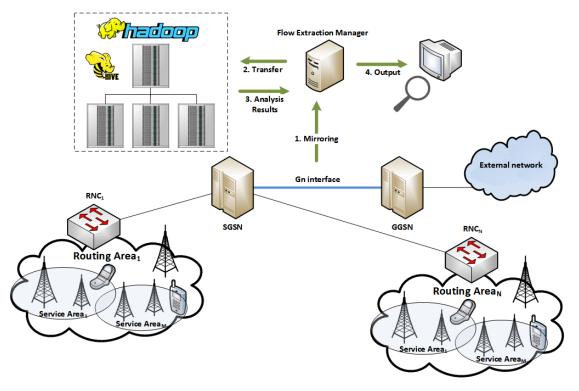


Figure 1: Flow Information Extraction Architecture

HSPA network, each subscriber will have one or more GTP tunnels, corresponding to each active PDP context, as well as possibly having separate tunnels for specific connections with different quality-of-service (QoS) requirements. Each tunnel is identified by a tunnel endpoint identifier (TEID) in the GTP-U messages, which should be a dynamically allocated random number. The control plane packets on the Gn interface are encapsulated on the control section of the GTP, namely GTP-C [9]. When a subscriber requests a PDP context, the SGSN will send a Create PDP Context Request GTP-C message to the GGSN giving details of the subscriber's request. The GGSN will then respond with a Create PDP Context Response GTP-C message which will either give details of the activated PDP context or will indicate a failure and give a reason for that failure. We make use of a Hadoop based platform (explained below) to aggregate packets in flows and correlate GTP-U and GTP-C information (described in Section 3.3). In this way we infer the mobility characteristics of the flows and associated GTP tunnels.

3.2.2. Mirroring of Gn interface

The network of interest in this paper consists of an area covered by 10 SGSNs. The average total traffic over all regional areas consists of approximately 15 billion packets in uplink direction and 20 billion packets in the downlink direction daily. This corresponds to approximately 80 TB of total data flowing in uplink and downlink daily in the mobile operator's core network. The importance of this work can be seen by the exponential increase in data traffic that has to be handled by a mobile operator. For example, in 2012, the approximate total data traffic was over 7TB in both uplink and downlink daily. The packets are captured by a mirroring device which was already in place on the operator premises as part of their already deployed measurement system. Hence, this work takes benefit of already deployed probes without requiring any additional deployment. We tested our method on real-world Gn interface Internet traffic data. The flow traces obtained from the mobile operator are collected by a server on a high speed link of 200 Mbit/sec at peak hours between 8-9 pm.

3.2.3. Hadoop platform

Among the available Big Data platforms, Hadoop [14] stands out as the most notable one as it is an open source solution. It is made up of a storage module, namely Hadoop Distributed File System (HDFS) and a computation module, namely MapReduce. Whereas HDFS can have centralized or distributed implementations, MapReduce inherently has a distributed structure that enables it to execute jobs in parallel on multiple nodes.

3.3. Flow Information and Mobility Characteristics extraction

The process proposed in this work, extracts and matches the user data TEID (TEID_DATA) field in the GTP-C messages (Create and Update PDP Context) with the TEID in GTP-U packets, in order to add the correspond-

ing location information to the traffic in the user plane.² GTP-C messages include an Information Element containing the location of the mobile terminal, expressed by Cell Global Identification (CGI), which is formed of the Mobile Country Code (MCC), the Mobile Network Code (MNC) the location area code (LAC), (which corresponds to the Routing Area identifier) and the cell identification ID (CELL-ID) or service area code (SAC) in Universal Mobile Telecommunications Service (UMTS). The service area (SA) identifies an area of one or more cells of the same location area and it is identified with a SAC, unique within that location area. As the MCC and MNC will remain unchanged for every packet in our trace, we monitor the LAC and SAC fields. A LA is a set of cells, which are grouped to decrease signalling overhead. Larger LAs reduce signalling for location updates because users hardly move out of the LA. However, the overhead introduced by paging is very high because there are many cells. Therefore, the size of location areas for circuit switched services can be larger than for packet switched services, which is why the term RA, which is a sub-division of a LA, is introduced in packet switched services. The differentiation between the RAs and LAs depends on the decision of the network operator. Typically, tens or even hundreds of base stations are present in a given LA. As we work with data transmissions, we will use the term RA from now on to refer to the greater area that groups several SAs.

A TEID uniquely identifies a tunnel endpoint on the receiving end of the GTP tunnel [9]. A local TEID value is assigned at the receiving end of a GTP tunnel in order to send messages through the tunnel. The GTP-C packets contain the information that identifies the location of the user (LAC and SAC fields) as well as a TEID_DATA field, which points to the identifier of the corresponding tunnel in the data plane. We extract from the GTP-U packets, which are identified by their TEID, several information regarding the characteristics of the mobile traffic, such as packet size or duration of the flows. The flow information can be extracted from data packets by taking into account simply five tuples namely, layer 4 protocol information (e.g. transport control protocol (TCP), User Datagram Protocol (UDP), Internet Control Message Protocol (ICMP), etc), source and destination IP address and source and destination port numbers. Taking advantage of their uniqueness and matching the values of TEID and TEID_DATA fields on both data and control planes respectively, the GTP-C and GTP-U information is joined by the Hadoop cluster through successive map and reduce operations, to obtain a merged table showing the flow and CELL-ID of the user flows. The complexity of the operation comes from the fact that, in a given PDP session with a specific TEID, there can be multiple flows (although each TEID belongs to a specific user). This must be taken into account while processing the data.

The result of the processing of the data mirrored from the Gn interface is a collection of anonymised rows, each representing a flow, providing information regarding the time it is collected, its size, a flow identifier, TEID and unique (SAC-LAC) information. This last tuple uniquely identifies the location of the user generating the flow. This information is obtained by filtering packet information and aggregating it into flows, later this aggregated information is processed through several map reduce operations to merge the information coming from the control and data plane. The mechanism is applicable to all mobile operators using the 3GPP standard Gn interface and can be put into practice immediately, since it does not require the change of the monitoring platform of the operator.

3.4. Mobile data traffic analysis

In this section, we analyse the main characteristics of our dataset, built from the information captured at the Gn interface between a SGSN and GGSN, as described in Section 3.1. We have collected two 1-hour traces at times with different traffic load in the network, close to the "peak" and "low" hours.

The Hadoop cluster in our platform was implemented based on Cloudera's Distribution Including Apache Hadoop (CDH4) [21] version on four nodes with Intel Xeon E5-2670 CPUs, 32 cores, 20 TB hard drive and 132 GB RAM.

Table 1 gathers the main characteristics of our dataset, after parsing, filtering and properly grouping the information available in the trace file. First of all, we present the total amount of data and control information that we have analyzed, expressing it in terms of flows and tunnels, which will be the main unit used in our mobility analysis. In light of our measurements, with more than a million flows³ spread over more than 80 thousand data tunnels, approximately 99% of flows and 98% of tunnels do not experience a handover. Note that the operator experience an overhead due to mobility management for these tunnels (mobility is provided by default in current operator networks), although no mobility is required by them. It also calls our attention that the number of control messages exchanged linked to data tunnels (identified by the TEID_DATA field) that are not exchanging any data packet. That translates into a considerable overhead in the already challenged operator's network to maintain context for data tunnels that are not being used for any data transmission. Data flows, although more numerous than tunnels, have shorter duration, which results in a lower average number of handovers per flow than per tunnel. Note that the total number of TEID handovers in Table 1 aggregates the number of handovers from all the tunnels analyzed.

 $^{^2\}mathrm{Note}$ that this location information is not carried in the data plane packets.

 $^{^{3}}$ Note that the "number of flows analyzed" in Table 1 refers to the number of flows that include useful information for our latter mobility analysis, not the total of flows in the network.

	Low hour	Peak hour
Total duration of the traffic trace	1 hour	1 hour
Total number of control packets	2722234	4294077
Total number of data packets	18978264	28582691
Total number of flows	1655550	2982064
Flows analyzed	1054566	1926062
Total number of control TEIDs in C-trace	762102	1259128
Total number of data TEIDs in C-trace	671078	1020178
Total number of TEIDs in U-trace	83510	154650
Total number of data TEIDs analyzed	60932	116529
Number of data TEIDs that experience handover	1464	2413
Total number of flow handovers	7402	18854
Total number of TEID handovers	2624	4750
TEIDs without handover in data trace	59468	114116
	(97.6%)	(97.9%)
Data flows without handover in data trace	1047380	1907759
	(99.3%)	(99%)

Table 1: Summary results from the data traces.

Figure 2 shows the number of handovers per tunnel for the low hour (Fig. 2a), and peak hour (Fig. 2b) performed in our dataset. We have included for comparison the number of tunnels and flows not moving in the network (tallest bars in x=0; note the histogram in logarithmic scale). In the peak hour, the number of handovers increase, even though the higher number of RA changes are not frequent.

With regard to mobility, we highlight the high percentage of data traffic in the network that does not experience a handover (99% of flows and 98.48% of tunnels in the medium load trace). Still, this traffic is provided with mobility support, just as any other in the network, which leads to an inefficient usage of resources in current network deployments, where mobility is granted to all the traffic being transmitted. Therefore, future network deployments could take into account these mobile data characteristics to improve scalability and free network resources.

First of all, we do not evaluate the changes of SA, for similarity to an initial deployment of DMM, which would be done by replacing packet gateway (P-GW) or serving gateway (S-GW) [22] (terms for GGSN and SGSN in LTE nomenclature). In addition, SAs are groups of cells, and they can overlap, which could lead us to misleading conclusions about user mobility. In any case, and just to confirm that the low mobility between routing areas is not due to the size of the area considered, we have measured the SA changes within a given RA. The number of tunnels that changes SA accounts to 5.6% of the tunnels analysed in our dataset, which confirms our previous conclusion and the low mobility in the traffic analysed.

Figure 3 represents the duration of tunnel identifiers and the median of the time between routing area handovers for every tunnel. As shown in Fig. 3 the points are concentrated in the lower X and Y range, meaning that from the flows experiencing a handover, there are a majority which are short lived and highly mobile. It can also be seen how the points form several lines across the plot (e.g., X = Y). These lines correspond to the points which ratio between the average time between handovers and TEID active time is constant. In fact these lines show the number of handovers that are more common among the users. The first of the lines (starting from the top of the graph) corresponds to tunnels surviving one handover. It is also worth explaining the reason behind the lack of data for X < Y. We measure the time between handovers through the changes in LAC and SAC of the active tunnels, hence for X < Y the tunnel is already disconnected and data is not available.

4. Understanding DMM deployment on current networks

In this section, we briefly explain how PMIPv6 and network-based DMM work and evaluate the usage of DMM as an alternative to current deployments and we use the mobile data traffic characteristics already presented to back the reasons behind our comparison with an example from actual data. As we have explained in Section 3.1, the tunneling protocol chosen to handle mobility in the UMTS network is GTP. However, given the increasing bandwidth demand, the characteristics of the mobile data traffic and the mobility requirements, we claim that mobility can be managed more efficiently.

4.1. PMIPv6

PMIPv6 is a network-based protocol that provides localized mobility management [23]. Its operation is conceptually identical to the one currently used in cellular networks, GTP. PMIPv6 releases the mobile terminal from its involvement in the mobility management of traditional IP mobility protocols, to place it at the network entities:

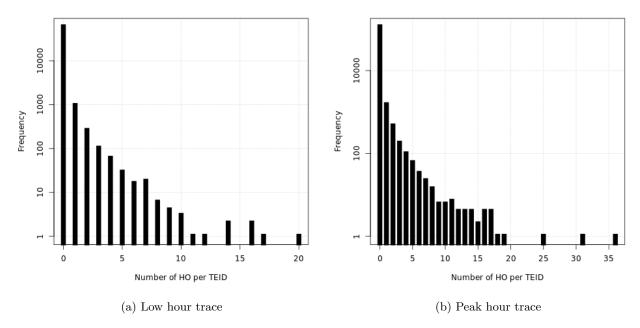


Figure 2: Histogram for the number of handovers performed per tunnel for the low and peak hours.

the Local Mobility Anchor (LMA) and the Mobile Access Gateway (MAG), as illustrated in Figure 4a. The LMA is placed at the operator's core network and manages the routing for all the mobile nodes within the mobility domain. Traffic for each mobile node is tunneled between the LMA and the corresponding MAG, through an IP in IP tunnel. The MAG performs all the mobility management signaling with the LMA on behalf of the mobile node by the exchange of Proxy Binding Update (PBU) and Proxy Binding Acknowledgment (PBA) messages. The MAG is the access router of the mobile node, that is, it acts as the first Internet Protocol (IP) hop in the mobility domain. In case of handover, the target MAG signals the LMA upon the attachment of the mobile node, establishing a new tunnel and eliminating the tunnel to the former MAG. The main problem of such a hierarchical architecture is the scalability, since all traffic must traverse the GGSN (LMA) which constitutes a bottleneck and single point of failure. In addition all traffic is treated the same and is provided with mobility management (is tunnelled from LMA to MAG).

4.2. Network-based DMM

DMM proposes a flatter architecture, placing the mobility management entities closer to the mobile node. DMM solutions are still being discussed at the IETF.⁴ For this work we consider our proposal in [24] as the basis for DMM operation, which is illustrated in Figure 4b. The roles of LMA and MAG are absorbed by a new entity, the Mobility Anchor and Access Router (MAAR). The MAAR

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acts as the Access Router for the mobile node, handles its mobility and the routing. Although there are different approaches, part of the functionality of the LMA is covered by a node that stores the prefixes assigned to every mobile node, mainly as a database (CMD in Figure 4a). In the case of handover, the target MAAR establishes a tunnel to the former MAAR to forward the traffic for the mobile node's previous prefix, while traffic not requiring mobility is forwarded directly to the Internet. Complexity increases as the mobile node performs more handovers, but in the general case DMM results in a more efficient and scalable approach for mobility management. One of the key aspects of DMM is that the flows originated by the mobile terminal are not tunneled, but only the flows performing handovers require the overhead of tunnelling through the network. For networks with low mobility characteristics, this is a key difference, since most of flows will not require differentiated handling from standard IP routing, reducing the overhead in the network.

4.3. DMM deployment analysis

Based on our previous analysis published in [25], we compare the packet delivery gains attained by the use of DMM principles. In a DMM architecture, the number of active prefixes (for relating to our scenario, active tunnels) is directly related to the number of handovers performed by the mobile node. In [25], it is proven that the gains of DMM over PMIPv6 are given by Equation 1, where \overline{N}_{PR} is the average number of active prefixes at a handover and $C_{MAAR-MAAR}$ and $C_{LMA-MAG}$ are expressed in terms of delay between MAARs and between LMA-MAG respec-

⁴http://datatracker.ietf.org/wg/dmm/documents/

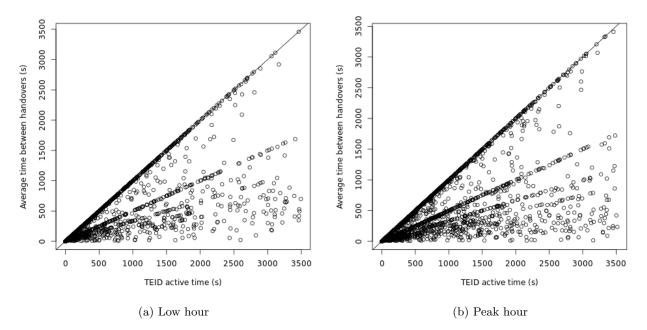


Figure 3: Duration of tunnel identifiers vs average of the time between handovers for the low and peak hours.

tively.

$$\frac{C_{DMM}}{C_{PMIP}} = (\overline{N}_{PR} - 1) \frac{C_{MAAR-MAAR}}{C_{LMA-MAG}}$$
(1)

The data on which we base our analysis comes from a Gn interface and corresponds to traffic between a GGSN and a single SGSN that handles several routing areas. We propose the scenario in Fig. 5 to compare with PMIPv6 and DMM approaches. We can see that PMIPv6 is conceptually similar to GTP, and the packet delivery cost comparison between DMM and PMIPv6 can be similarly extended to GTP. We propose as reasonable scenario to have the LMA co-located with the GGSN and the MAG with the SGSN. For the DMM architecture, we propose to place a MAAR in every RA (i.e. RNC). Therefore, to follow the analysis given by Equation 1 the number of active prefixes applying DMM to our scenario is given by the number of tunnels performing a RA change. As the analysis in [25] involves the average number of prefixes active per user and we have the aggregated information of active tunnels in the network, for comparison with PMIPv6, we need to take into account all the active tunnels in the network (which is equal to 1 active prefix per user in the analysis provided in [25]). We have tracked the RA changes per tunnel, as reported in Section 3.4, to find out that the average number of handovers per TEID in the network is 1.8 and the total number of tunnels performing at least one handover is 1464 for the low hour and 2413 for the peak hour (roughly 2% of the tunnels in the network for both cases). With this information and the number of packets delivered per tunnel, we can say that the packet delivery cost ratio is given by Equation 2, where λ_{mov} is the number of data packets sent in the tunnels that experienced

handover, λ_{tot} is the total number of data packets, N_{mov} is the number of tunnels that experience handover and N_{tot} is the total number of tunnels. This is equivalent to the ratio between the traffic that experiences a handover (and should be provided mobility) and the total amount of traffic being served in the network. Note that, for λ_{mov} we are considering a worst case scenario, as the amount of data packets that would be tunnelled by DMM would be the packets that are sent after the handover, but not in the routing area in which the flow is originated.

$$\frac{C_{DMM}}{C_{PMIP}} = \frac{\lambda_{mov}/N_{mov}}{\lambda_{tot}/N_{tot}} \frac{D_{MAAR-MAAR}}{D_{LMA-MAG}}$$
(2)

Table 2: Value of the parameters in Eq. 2

	Low hour	Peak hour
λ_{mov}	41467	49927
λ_{tot}	18978264	28582691
N_{mov}	1464	2413
N_{tot}	60932	116529

The values for the two datasets analyzed to be substituted in Equation 2 are shown in Table 2. Applying these results to Equation 2 for the peak hour trace we obtain:

$$\frac{C_{DMM}}{C_{PMIP}} = 0.08 \frac{D_{MAAR-MAAR}}{D_{LMA-MAG}} \tag{3}$$

Computing the delay between MAARs that makes equal the packet delivery cost of applying PMIPv6 and DMM, we have for the low and peak hours respectively:

$$D_{MAAR-MAAR}^{low_hour} = 11.85 D_{LMA-MAG} \tag{4}$$

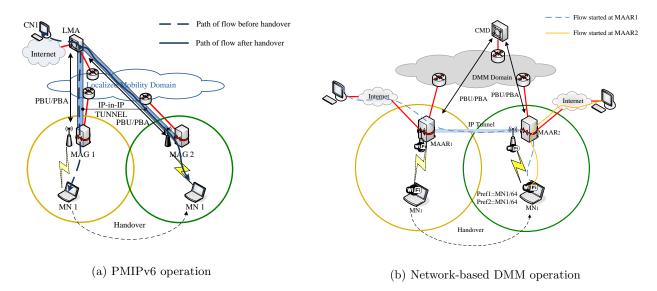


Figure 4: Operation of PMIPv6 and network-based DMM protocols.

$$D_{MAAR-MAAR}^{peak_hour} = 10.99 D_{LMA-MAG} \tag{5}$$

This result can be explained given the low mobility in the network. For instance, in the peak hour case, the number of tunnels that are provided mobility is so low compared to the total amount of traffic in the network that the distance between MAARs (in terms of delay) can be up to 11.85 times the distance between LMA and MAG for the packet delivery cost of DMM and PMIPv6 to be equal. This analysis, based on real traces from operator's network, indicates that this network can be flattened by a factor of 11, with no extra overhead in the network, reducing the use of resources in the core of the operator and providing on-demand mobility to the user. We claim the use of this technology is key for the development of 5G network architectures, characterized by extreme consumption of resources in the network. The mobility in the network for the two loads studied is around 2%, despite the amount of traffic served in the system. This favours the deployment of a flatter architecture like DMM as the traffic that requires mobility is a small portion of the total. In addition, TEIDs, in average do not experience a high number of RA changes, so the number of tunnels that should be maintained active in a DMM solution would be much lower.

5. Conclusion

In this work, we collect the information observed in the Gn interface between the core network nodes and extract handover-related information for further analysis from a real operator's network. The extraction steps, which can be applied to all mobile networks using 3GPP standard, are provided in detail. We match the location information present in the control plane, which is missing in the user plane packets to evaluate mobility in the network thanks to

big data platform and analyze the mobility requirements. Our mobility analysis reveals that a DMM-based mobility solution, flattening the current deployment by a factor of almost 11, can be applied. The new mobility management solution will provide increased scalability to the operator network, without any additional overhead. The limits on the profitability of this solution will depend on the specific characteristics of the operator's network and the mobile data traffic mobility. As future work we plan to extend the period of time analyzed, searching for time patterns and providing guidelines for deployment of mobility solutions.

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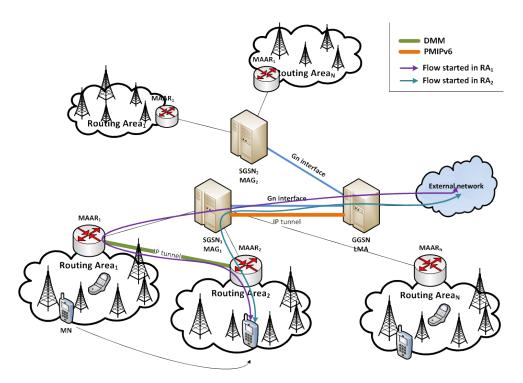


Figure 5: Cellular network architecture with PMIPv6 and DMM mobility entities co-located.

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