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RESEARCH ARTICLE

Providing Throughput Guarantees in Heterogeneous Wireless Mesh Networks

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ABSTRACT

In this paper we propose to provide throughput guarantees in heterogeneous wireless mesh networks by jointly optimizing routing and MAC configuration. Our solution is based on the notion of *linearized capacity region*, which provides a technology-independent way of representing the capacity of a wireless link (thereby hiding the technology specifics to the upper layers). Based on the available capacity of the underlying links as given by the linearized capacity region, we propose two routing algorithms (based on multi-path and single-path, respectively) that find optimal paths for all the flows in the network given their throughput requirements. The throughput allocation resulting from routing is then provided to each link, which uses this information to optimize its technology-specific MAC parameters. The proposed approach is evaluated in an heterogeneous scenario comprising WLAN and WiMAX technologies, and is shown to outperform previous solutions by (at least) a factor of 2.

KEYWORDS

Wireless, IEEE 802.11, Capacity Region

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1. INTRODUCTION

In the last few years, mesh networking has emerged as a cost effective and efficient solution for realizing backhaul networks. The multihop wireless network architecture of mesh networks enables them to efficiently cover large areas without requiring many interconnections into a wired infrastructure. Furthermore, mesh networks are dynamically self-organized and self-configured, which ultimately results in reduced up-front cost and lower network maintenance costs for the operator. Along these lines, many major operators have already considered wireless mesh networks (WMNs) as a technology for their wireless Cities initiatives [1][1].

Most existing WMN solutions are designed for a single, specific radio technology. However, combining different technologies to realize a heterogeneous mesh solution allows for a more flexible designs that benefit from the complementary characteristics of different technologies, e.g., extending the wireless mesh connectivity of WLAN access points, using a point-to-multipoint non-line-of-sight technology like WiMAX.

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A critical concern for operators is to provide their customers with service guarantees. While already challenging in WMNs with single radio technology, this becomes even more difficult in heterogeneous mesh networks. This is because different radio technologies typically exhibit vastly different link characteristics, in particular in the way capacity is shared between links, and this needs to be considered in admission control, routing, traffic engineering, etc.

In this paper, we aim at designing a solution to provide *service guarantees in heterogeneous WMNs*, while preserving flexibility and cost-efficiency through a *technology-independent approach*. The paper first analyzes the capacity region of each wireless link and then proposes a routing algorithm that optimizes performance within the capacity region of all wireless links of the mesh. Our key contributions are: i) We propose a novel technology-independent way to represent the capacity region of a wireless link, hereafter referred to as linearized capacity region. ii) We provide a method to map the capacity region of a 802.11 link to the proposed linearized capacity region. iii) We design a novel routing algorithm for heterogeneous WMNs that builds upon the linearized capacity region concept. We evaluate the performance of the resulting approach, by means of simulations for heterogeneous WMNs comprising 802.11 and 802.16 technologies, and show that it outperforms previous approaches by at least a factor of 2.

The rest of the paper is structured as follows. In Section 2 we propose the linearized capacity region concept, which is mapped to the 802.11 technology in Section 3. Based on this concept, in Section 4 we present a routing algorithm for heterogeneous WMNs. This algorithm is evaluated and compared against other routing approaches in Section 5. Finally, Section 6 reviews some related work and Section 7 closes the paper with some final remarks.

2. LINEARIZED CAPACITY REGION

In order to allocate resources in a network that comprises wireless links, one needs to know the capacity region of those links. However, these regions are typically very complex to compute and depend on the specific technology used for each wireless link. In this paper we advocate for the need of a simple and technology-independent way of representing the capacity region of a wireless link, in order to reduce the complexity of resource allocation algorithms[2].

Following the above, in this section we present a novel approach to represent the feasible allocations in a wireless link. The key advantages of the proposed approach are twofold: i) The approach is technology-independent and can be mapped to different technologies as we do in Section 3. Thus, we can use it for resource allocation in heterogenous mesh networks. ii) The proposed solution can be easily combined with *efficient* algorithms to optimize the mesh network performance, as we do in Section 4.

2.1. Capacity region of a wireless link

While the available capacity of a wired link is a precise notion, it is a blurred notion for wireless links. Take for instance a WLAN link with two nodes. If all the link capacity is allocated to one of the nodes, the available capacity equals the WLAN nominal rate[3] as no time is wasted in collisions. However, if capacity is fairly shared between the two nodes, the available capacity is smaller due to the bandwidth wasted in contention.

We note that, while it may be theoretically possible to compute the exact capacity region of all the technologies and links of an heterogeneous WMN, designing an algorithm for resource allocation that relies on such complex and technology dependent computations would be highly inefficient. Instead, in the following we propose a novel notion to characterize the capacity region of a wireless link in a simple and technology-independent way. This notion allows to easily determine the set of resource allocations that are feasible in the link.

2.2. Proposed concept

Let us consider a wireless link L shared by N flows, where a flow corresponds to the traffic from a given source node to another destination node in the link. Let R_i denote the throughput allocated to flow i. The key problem of

[2] The need for a technology-independent representation the available capacity of a link has already been detected by the 802.21 standard [5], which includes a primitive for this purpose. However, the 802.21 does not address the actual meaning of available capacity in wireless systems, and leaves the computation and interpretation of this parameter up to the implementation.

[3] By the *nominal rate* of a technology, we mean the highest data rate provided by this technology. For instance, for the 802.11b technology we mean a data rate of 11 Mpbs.

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allocating resources in link L is to determine whether a given allocation $\{R_1, \ldots, R_N\}$ of flows in this link is feasible or not. In wired links this is straightforward: as long as the total resources in the link do not exceed the link's capacity C, the allocation is feasible, and otherwise it is not feasible, i.e.

$$\sum_{i \in L} R_i \leq C \Longleftrightarrow \{R_1, \ldots, R_N\} \text{ is feasible} \quad (1)$$

Determining the feasible allocations in a wireless link is much more difficult because, in contrast to wired links, the total amount of resources allowed is not constant but depends on a number of factors including, e.g. the wireless technology used in the link (contention-based technologies waste some resources in collisions which centralized technologies do not) or the modulation and coding scheme used by each of the nodes of the link[4].

In order to provide a way of expressing the capacity region of the wireless link as accurately as possible while avoiding the complexity involved in considering all the above aspects, our key proposal is to *linearize the capacity region of a wireless link*. In particular, with our proposal any allocation that satisfies

$$\sum_{i \in L} c_i R_i \leq C \quad (2)$$

is guaranteed to be feasible, where c_i is defined as the cost of flow i and C as the wireless link capacity. Hereafter, we refer to the capacity region resulting from this linearization as the *linearized capacity region*. Note that this way the problem of allocating resources in a network with wireless links becomes as easy as with wired links, as we can determine it with a simple linear function. This concept is shown in Figure 1.

The shortcoming of this approach is that since we are using a lower bound of the actual capacity region, there may be some feasible zones in this region that are not allowed, and this may lead to a suboptimal allocation that does not take full advantage of the wireless link feasible allocations. However, the results presented in Section 5 show that there is no significant performance loss because of this reason.

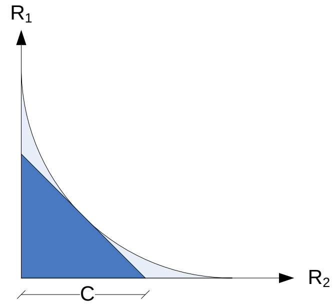

Figure 1. Linearized capacity region.

2.3. Mapping to wireless technologies

The challenge of the proposed model is the computation of the c_i and C values for the different wireless technologies. For technologies with centrally coordinated medium access such as IEEE 802.16, this computation is rather direct. Indeed, as centralized approaches do not waste any bandwidth in contention, channel time is fully used and as a result the total capacity depends only on how time is shared among the different stations and the modulation rate that these are using. Therefore in this case the boundary of the capacity region can be computed as

$$\sum_{i \in L} c_i R_i = C \quad (3)$$

where C is the nominal bit rate of the technology and c_i is the ratio between this nominal bit rate and the rate of the modulation scheme used by flow i[5].

The obtention of the linearized capacity region parameters for contention based wireless technologies is more difficult. The following section is devoted to the computation of these parameters for IEEE 802.11[6].

3. LINEARIZED CAPACITY REGION MAPPING TO IEEE 802.11

In this section we first present a model for 802.11 that, given the configuration of the contention parameters of the nodes of a link, computes the throughput obtained

[4] Hereafter we refer with *modulation rate* to the date rate provided by the modulation and coding scheme used. For example, the modulation rates available with 802.11b are 11 Mbps, 5.5 Mbps, 2 Mbps and 1 Mbps.

[5] The reader is referred to Section 5.5 for the specific C and c_i values of a typical IEEE 802.16 configuration.

[6] The mapping to contention based technologies other than IEEE 802.11 can be performed following a similar method to the one outlined in the following section.

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by each node. Then, we propose a bandwidth allocation strategy that, given a desired throughput allocation, finds (if it exists) the optimal configuration that satisfies this allocation. Finally, we compute the linearized capacity region resulting from our throughput allocation strategy.

3.1. Bandwidth allocation model

The resulting bandwidth allocation in 802.11 depends on the CW_{min} and CW_{max} parameters of each node, which in the latest release of the standard are configurable parameters [6]. Following our previous results of [7], in this paper we take $CW_{min} = CW_{max}$, as [7] shows (both analytically and via simulation) that no other configuration provides better throughput performance. We denote by CW_j the CW configuration of node j.

Following the analysis of [8], if the backoff window size of a station is constant, the probability that it transmits a packet in a slot time is given by[7]

$$\frac{2}{CW_j + 1} \tag{4}$$

Let $\rho_{i,j}$ denote the probability that a packet transmitted by station j belongs to flow i. Let us further denote by τ_i the probability that flow i makes a transmission attempt in a randomly chosen slot time. Then,

$$\tau_i = \rho_{i,j} \frac{2}{CW_j + 1} \tag{5}$$

The throughput provided to a flow can be computed as a function of the τ_i's as follows [8][8]:

$$r_i = \frac{p(s_i) l}{p(s)T_s + p(c)T_c + p(e)T_e} \tag{6}$$

where l is the average packet length, T_s, T_e and T_c are the average durations of a successful slot time, an empty one and a collision, respectively, and $p(s_i)$, $p(s)$, $p(c)$ and $p(e)$ are the probabilities that a slot time contains a success of flow i, a success of any flow, a collision and is empty, respectively.

The probabilities are computed as

$$p(s_i) = \tau_i \left(\prod_{j \in S \setminus S_i} \left(1 - \sum_{k \in F_j} \tau_k \right) \right) \tag{7}$$

$$p(s) = \sum_{i \in F} p(s_i) \tag{8}$$

$$p(e) = \prod_{j \in S} \left(1 - \sum_{k \in F_j} \tau_k \right) \tag{9}$$

$$p(c) = 1 - p(e) - p(s) \tag{10}$$

where S denotes the set of stations in the link, S_i is the station to which flow i belongs, F is the set of flows in the link, and F_j is the set of flows of station j.

We can use the following approximations for the above probabilities[9]:

$$p(s_i) = \tau_i \left(1 - \sum_{j \in S \setminus S_i} \sum_{k \in F_j} \tau_k \right) \tag{11}$$

$$p(e) = 1 - \sum_{i \in F} \tau_i + \sum_{i \in F} \sum_{j \in F \setminus \{F_1, \ldots, F_{S_i}\}} \tau_i \tau_j \tag{12}$$

$$p(c) = \sum_{i \in F} \sum_{j \in F \setminus \{F_1, \ldots, F_{S_i}\}} \tau_i \tau_j \tag{13}$$

The average slot time durations are computed as

$$T_s = \frac{1}{p(s)} \sum_{i \in F} p(s_i) T_{s,i} \tag{14}$$

$$T_c = \frac{1}{p(c)} \sum_{i \in F} \sum_{j \in F \setminus F\{F_1, \ldots, F_{S_i}\}} \tau_j \tau_i T_{c,i,j} \tag{15}$$

where $T_{s,i}$ is the duration of a successful transmission of flow i and $T_{c,i,j}$ is the duration of a collision between flows i and j.

$T_{s,i}$ is computed as follows

$$\begin{aligned} T_{s,i} = T_{PLCP} + \frac{H + l}{C_i} &+ SIFS + T_{PLCP} \\ &+ \frac{ACK}{C_i} + DIFS \end{aligned} \tag{16}$$

[7] Note that our definitions of slot time duration and probability follow the notion of a slot time given in [8], which defines a slot time as the period elapsed between two backoff counter decrements.

[8] We note that, like in [8], our throughput analysis assumes saturation conditions, i.e. stations always have packets ready for transmission. We argue that this assumption is appropriate to compute the bandwidth allocated to a flow, since it provides the maximum throughput that can be obtained by the flow.

[9] Note that these are worst-case approximations that lead to smaller throughputs than the actual ones. This is important because it ensures that the result of linearizing the approximate capacity region yields a feasible region.

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where T_{PLCP} is the Physical Layer Convergence Protocol preamble and header transmission time, H is the MAC overhead (header and FCS), ACK is the size of the acknowledgment frame, l is the packet length, $SIFS$ and $DIFS$ are time constants defined by the standard [6], and C_i is the bit rate of the modulation scheme used for flow i.

Finally, $T_{c,i,j}$ is computed as

$$T_{c,i,j} = T_{PLCP} + \frac{H+l}{min(C_i, C_j)} + EIFS \tag{17}$$

where $EIFS$ is another time constant defined by the standard.

3.2. Bandwidth allocation strategy

Based on the model presented above, in the following we address the problem of finding the optimal 802.11 configuration to meet a given set of bandwidth requirements. In particular, given a set of desired throughputs for each flow in the link, $\{R_1, \ldots, R_N\}$, our goal is to find the set of CW_j's that meets these throughput requirements, i.e. the configuration that provides each flow i with an allocated throughput r_i no smaller than its desired throughput:

$$r_i \geq R_i \tag{18}$$

The first step towards finding the CW_j configuration is to find the τ_i values that provide each flow with the desired throughput. These τ_i values are found with the algorithm described next. We first distribute throughput among the competing flows proportionally to their allocated rate,

$$\frac{r_i}{r_j} = \frac{R_i}{R_j} \tag{19}$$

which, combined with Eq. (6), yields

$$\frac{\tau_i(1-\sum_{k\in S_j}\tau_k)}{\tau_j(1-\sum_{k\in S_i}\tau_k)} = \frac{R_i}{R_j} \tag{20}$$

The above can be approximated by $\frac{\tau_i}{\tau_j} \approx \frac{R_i}{R_j}$.

Based on the above equation, we can express all τ_i's as a function of a reference τ_1 as follows:

$$\tau_i = w_i \tau_1 \tag{21}$$

where $w_i = R_i/R_1$.

With the above, the total throughput in the WLAN r can be computed as a function of τ_1 as

$$r = \sum_{i\in F} r_i = \frac{l\left(a\tau_1 + b\tau_1^2\right)}{c + d\tau_1 + e\tau_1^2} \tag{22}$$

where

$$\begin{aligned}
a &= \sum_{i\in F} w_i \\
b &= -\sum_{i\in F}\sum_{j\in F\setminus F_{S_i}} w_i w_j \\
c &= T_e \\
d &= \sum_{i\in F} w_i T_{s,i} - \sum_{i\in F} w_i T_e \\
e &= -\sum_{i\in F}\sum_{j\in F\setminus F_{S_i}} w_i w_j T_{s,i} \\
&+ \sum_{i\in F}\sum_{j\in F\setminus\{F_1,\ldots,F_{S_i}\}} w_i w_j T_{c,i,j} \\
&+ T_e \sum_{i\in F}\sum_{j\in F\setminus\{F_1,\ldots,F_{S_i}\}} w_i w_j
\end{aligned} \tag{23}$$

In order to find the optimal configuration, we look for the τ_1 value that maximizes the r_i's. Considering that throughput is distributed among flows following Eq. (19), this τ_1 value can be found by maximizing r; hence,

$$\frac{\partial r}{\partial \tau_1} = 0 \tag{24}$$

which yields

$$(a + 2b\tau_1)(c + d\tau_1 + e\tau_1^2) - (a\tau_1 + b\tau_1^2)(d + 2e\tau_1) = 0 \tag{25}$$

The above can be expressed as

$$A\tau_1^2 + B\tau_1 + C = 0 \tag{26}$$

where

$$\begin{aligned}
A &= bd - ea \\
B &= 2bc = -2\sum_{i\in F}\sum_{j\in F\setminus F_{S_i}} w_j w_i T_e \\
C &= ca = T_e \sum_{i\in F} w_i
\end{aligned} \tag{27}$$

From solving the above second order equation, we can isolate τ_1 and then compute the τ_i's from Eq. (20).Once the τ_i values have been obtained, we can compute the CW_j

values as follows. From Eq. (5),

$$\sum_{i\in F_j} \tau_i = \sum_{i\in F_j} \rho_{i,j} \frac{2}{CW_j+1} = \frac{2}{CW_j+1} \tag{28}$$

Finally, by isolating CW_j from the above we obtain

$$CW_j = \frac{2}{\sum_{i\in F_j} \tau_i} - 1 \tag{29}$$

which terminates the configuration of the CW parameters for bandwidth allocation.

3.3. Capacity region computation

In the following we compute the *linearized capacity region* of a WLAN link that is using the strategy described above to allocate bandwidth to its flows. In particular, we compute the costs c_i and the wireless link capacity C of Eq. (2), which provides a mapping of the *linearized capacity* concept to the 802.11 wireless technology. In order to linearize the capacity region of a WLAN, we first have to choose the tangent point of the linearized region. Indeed, we can choose different tangent points $\{R_1, \ldots, R_N\}$ which result in different linearized regions. Note that if we chose a tangent point in which $R_i > R_j$, we are favoring flow i over flow j, as the linearized capacity region is closer to the actual capacity region for those allocations where flow i takes larger throughputs.

Fig. 2 illustrates three different tangent points that can be chosen to build the linearized region. Tangent point (b), for which $R_1 > R_2$, favors flow 1 over flow 2 as it covers well the area where the throughput of flow 1 is high while it leaves out the area where flow 2 has a large throughput. In contrast, tangent point (c) favors flow 2 as it covers well the area where $R_2 > R_1$.

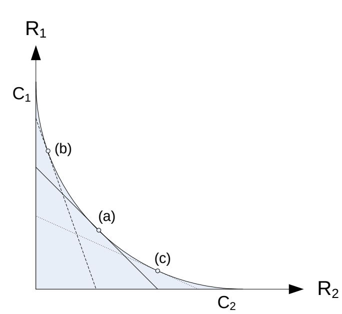

Figure 2. Tangent point options.

In order to choose the tangent point for the linearized capacity region, we would like to find an appropriate compromise between favoring those flows that are using higher modulation rates, as allocating throughput to these flows yields a more efficient utilization of the overall wireless resources and starving those flows which are using lower modulation rates, since in case one of these flows is needed (for instance, because it belongs to a critical path for routing), we would like to be able to use it.

The proportional fairness criterion has been defined precisely to satisfy the above compromise. In this paper we take the *proportional fair* allocation as the tangent point for the linearized capacity region. In particular, the proportional fair allocation is the one that satisfies

$$\frac{R_i}{R_j} = \frac{C_i}{C_j} \tag{30}$$

where C_i and C_j are the modulation rates of flows i and j, respectively [9].

Note that in the graph of Fig. 2, the proportional fair allocation corresponds to tangent point (a) which is indeed a compromise between allocating more throughput to the better off flows while not starving the worst off ones (note that the graph uses different scales for the axes).

The above tangent point can be easily obtained by using the model of Section 3.2 together with Eq. (30). We take the following approach.

According to Eqs. (2) and (30), our goal is to find the c_i and C parameters such that

1. Any allocation $\{R_1, \ldots, R_N\}$ that satisfies $\sum_i c_i R_i \le C$ falls within the capacity region.
2. The proportional fair allocation satisfies equality, i.e. $\sum_i c_i R_i = C$.

Let us consider the function $\sum_i c_i r_i$, where $\{r_1, \ldots, r_N\}$ are the boundaries of the capacity region as computed in Section 3.2. With this function, we proceed as follows to compute the parameters of the linearized capacity region:

$i)$ We first compute the parameters c_i by imposing that $\sum_i c_i r_i$ is minimized for the $\{r_1, \ldots, r_N\}$ point that corresponds to the proportional fair allocation.

$ii)$ Next, we obtain the value of C by evaluating the function $\sum_i c_i r_i$ at this point.

We next show that with this procedure we satisfy the two objectives stated above. Indeed, objective 2) is clearly

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satisfied by $ii)$. On the other hand, $i)$ imposes that all the points in the boundary of the capacity region satisfy $\sum_i c_i r_i \geq C$. This implies that any point that satisfies $\sum_i c_i R_i \leq C$ will fall within the capacity region, and therefore objective 1) is also met.

Following the above, the remaining challenge is to compute the costs c_i for which function $\sum_i c_i r_i$ takes a minimum at $R_i = w_i R_1$, where flow 1 is taken as reference and

$$w_i = \frac{C_i}{C_1},\ i \in \{2, \ldots, N\} \tag{31}$$

The above yields the following system of equations:

$$\left.\frac{\partial \sum_j c_j r_j}{\partial w_i}\right|_{w_i = C_i/C_1} = 0,\ i \in \{2, \ldots, N\} \tag{32}$$

Substituting r_i by the expression given in Eq. (6) and operating on the equation, we obtain:

$$\begin{aligned} & \frac{\partial \sum_j c_j p(s_j)}{\partial w_i} \left(p(s)T_s + p(c)T_c + p(e)T_e\right) \\ - & \sum_j c_j p(s_j) \left(\frac{\partial p(s)T_s}{\partial w_i} + \frac{\partial p(c)T_c}{\partial w_i} + \frac{\partial p(e)T_e}{\partial w_i} \right) = 0 \end{aligned} \tag{33}$$

We compute the first term of the above equation as follows

$$\frac{\partial \sum_j c_j p(s_j)}{\partial w_i} = c_i \frac{\partial p(s_i)}{\partial w_i} + \sum_{j \in F \setminus i} c_j \frac{\partial p(s_j)}{\partial w_i} \tag{34}$$

where

$$\begin{aligned} \frac{\partial p(s_i)}{\partial w_i} = \tau_1 + w_i \frac{\partial \tau_1}{\partial w_i} & - \sum_{j \in F \setminus F_i} w_j \tau_1^2 \\ & - \sum_{j \in F \setminus F_i} w_i w_j 2\tau_1 \frac{\partial \tau_1}{\partial w_i} \end{aligned} \tag{35}$$

and

$$\frac{\partial p(s_j)}{\partial w_i} = w_j \frac{\partial \tau_1}{\partial w_i} - w_j \tau_1^2 - \sum_{k \in F \setminus F_j} w_k w_j 2\tau_1 \frac{\partial \tau_1}{\partial w_i} \tag{36}$$

for $F_i \neq F_j$, and

$$\frac{\partial p(s_j)}{\partial w_i} = w_j \frac{\partial \tau_1}{\partial w_i} - \sum_{k \in F \setminus F_j} w_k w_j 2\tau_1 \frac{\partial \tau_1}{\partial w_i} \tag{37}$$

for $F_i = F_j$.

To compute $\partial \tau_1 / \partial w_i$, we take the expression of τ_1 given by Eq. (26) and proceed as follows:

$$\begin{aligned} \frac{\partial \tau_1}{\partial w_i} & = \frac{\left(-B' + \frac{2BB' - 4(A'C + AC')}{2\sqrt{B^2 - 4AC}}\right) 2A}{4A^2} \\ & - \frac{2A'(-B + \sqrt{B^2 - 4AC})}{4A^2} \end{aligned} \tag{38}$$

where A', B' and C' are the derivatives of the parameters A, B and C given in Eq. (26):

$$\begin{aligned} A' & = \frac{\partial b}{\partial w_i} d + b \frac{\partial d}{\partial w_i} - \frac{\partial e}{\partial w_i} a - e \frac{\partial a}{\partial w_i} \\ B' & = -4T_e \sum_{j \in F \setminus F_i} w_j \\ C' & = T_e \end{aligned} \tag{39}$$

The partial derivatives included in the above expressions are computed as follows:

$$\frac{\partial a}{\partial w_i} = 1 \tag{40}$$

$$\frac{\partial b}{\partial w_i} = -2 \sum_{j \in F \setminus S_i} w_j \tag{41}$$

$$\frac{\partial c}{\partial w_i} = 0 \tag{42}$$

$$\frac{\partial d}{\partial w_i} = T_{s,i} - T_e \tag{43}$$

$$\begin{aligned} \frac{\partial e}{\partial w_i} & = - \sum_{j \in F \setminus F_i} w_j T_{s,i} - \sum_{j \in F \setminus F_i} w_j T_{s,j} \\ & + \sum_{j \in F \setminus F_i} w_j T_e + \sum_{j \in F \setminus F_i} w_j T_{c,i,j} \end{aligned} \tag{44}$$

Finally, the rest of the terms of Eq. (33) are computed as follows:

$$\frac{\partial p(s)T_s}{\partial w_i} = \sum_{j \in F \setminus i} \frac{\partial p(s_j)}{\partial w_i} T_{s,j} + \frac{\partial p(s_i)}{\partial w_i} T_{s,i} \tag{45}$$

$$\begin{aligned} \frac{\partial p(c)T_c}{\partial w_i} & = \sum_{k \in F} \sum_{j \in F \setminus F \in (F_1..F_k)} w_k w_j T_{c,k,j} 2\tau_1 \frac{\partial \tau_1}{\partial w_i} \\ & + \sum_{j \in F \setminus F_i} w_j T_{c,i,j} \tau_1^2 \end{aligned} \tag{46}$$

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$$\frac{\partial p(e)}{\partial w_i} = - \tau_1 - \sum_{j\in F} w_j \frac{\partial \tau_1}{\partial w_i} + \sum_{j\in F\setminus F_i} w_j \tau_1^2 + \sum_{k\in F} \sum_{j\in F\setminus F\in(F_1..F_k)} w_k w_j 2\tau_1 \frac{\partial \tau_1}{\partial w_i} \quad (47)$$

By substituting Eqs. (33)-(47) into Eq. (32), we obtain a system of $N-1$ equations on the c_i's. Note from Eq. (2) that we have one degree of freedom when fixing the c_i's and C. Therefore we can set without loss of generality $c_1 = 1$ and, as a result, we have a system of $N-1$ equations with $N-1$ unknowns. As this system of equations is linear, it can be easily solved by means of standard techniques like e.g. Cramer.

From the above, we can compute all the c_i's by resolving the system of equations and then we can compute C from

$$C = \sum_i c_i r_i\big|_{w_i=C_i/C_1} \quad (48)$$

which terminates the computation of the mapping to the the linearized capacity region.

4. ROUTING AND BANDWIDTH ALLOCATION

The key objective of the proposed linearized capacity region is to aid the design of efficient algorithms to optimize network performance. In particular, the proposed model allows to easily determine if a given flow allocation is feasible and therefore it is very useful for the design of efficient optimization algorithms. To illustrate this, in this section we present a routing algorithm for mesh networks that relies on the proposed linearized capacity region in order to provide throughput guarantees[10].

The specific optimization problem that we address in this section is stated as follows: Given a mesh network consisting of a set of wireless links with corresponding linearized capacity region, a set of gateways and a set of flows, with each flow i originating at source node N_i and having a throughput requirement R_i, we want to find a route for each flow to *any* of the gateways of the mesh network such that the throughput requirement of each flow is met and the number of admitted flows is maximized.

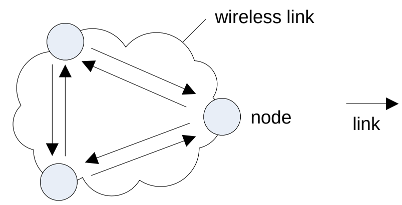

Figure 3. Wireless link with 3 nodes.

We assume that the mesh network has a proper radio resource management, such that packet transmissions do not fail because of interference from neighboring links. This assumption is supported by our measurements of [10], which show that as long as channel separation is large enough, different channels do not interfere with each other. Moreover, the use of rate adaptation techniques further mitigates the impact of interference.

We consider two scenarios: i) A multipath scenario in which each flow can be split and distributed over different paths and ii) A single-path scenario in which each flow is treated as an atomic, unsplittable entity and therefore it can be routed through only one single path.

Let us start with the multipath routing problem. This problem can be viewed as a *multi-commodity flow problem* [11], which has been widely studied in the literature, and can thus be formulated as the following linear programming (LP) problem.

Let $r_{i,l}$ be the rate allocated to flow i on link l, where here a *link* denotes a pair of directly connected nodes, as opposed to the notion of *wireless link* that we used in the previous sections to refer to a set of nodes that can communicate with each other by sharing a common wireless capacity (see Fig. 3 for an example of a wireless link with three nodes and the corresponding links). Let s_i denote the set of links originating at the source node of flow i. Furthermore, let N_{in} denote the set of links that terminate at node N and N_{out} the set of links that leave from this node. We further denote by I the set of nodes that are neither sources nor gateways. Following the above definitions, our objective is to find the allocation that satisfies

$$\min \sum_{i,l} r_{i,l} \quad (49)$$

[10] It is important to note that the presented algorithm is only an *example* to show the potential of the proposed concept. Indeed, the linearized capacity region can be used to solve other optimization problems such as, e.g., network planning, traffic engineering or admission control.

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subject to

$$\sum_{l \in s_i} r_{i,l} = R_i, \; \forall i \tag{50}$$

$$\sum_{l \in N_{in}} r_{i,l} = \sum_{l \in N_{out}} r_{i,l}, \; \forall i, N \in I \tag{51}$$

$$\sum_{l \in L} c_l \sum_{i \in l} r_{i,l} \leq C_L, \; \forall L \tag{52}$$

The expression of Eq. (49) aims at finding, among all solutions that satisfy the throughput requirements of all flows, the one that minimizes the sum of all individual rates. The reason behind this is to try to minimize, as long as the throughput requirements are met, the total number of hops of all flows in the mesh. Furthermore, Eq. (50) imposes that the sum of rates of a flow leaving the flow's source node is equal to the guaranteed throughput for this flow, which ensures that the desired throughput guarantees are satisfied[11]. Additionally Eq. (51) imposes the flow conservation constraints by guaranteeing that the sum of incoming rates to every node equals the sum of outgoing rates; note that this equation applies to all nodes but sources and gateways. Finally, Eq. (52) imposes the capacity constraints for each wireless link L as given by our linearized capacity model. In particular, this equation imposes that the sum of the aggregated rates for each link, weighted by the cost, cannot exceed the wireless link capacity C_L.

The above LP problem can be solved by using standard techniques. As a result of this, we obtain a routing strategy that admits as many flows as possible while meeting the desired throughput guarantees.

We next address the single-path routing problem. This can be viewed as another very widely studied problem in the literature which is the *unsplittable flow problem*, and can be formulated as the following integer programming (IP) problem.

Let x_i be 1 if flow i is routed and 0 otherwise. Furthermore, let $y_{i,l}$ be 1 if the path chosen for flow i traverses link l and 0 otherwise. Then, we want to find the allocation that satisfies

$$\max \sum_i x_i \tag{53}$$

subject to

$$x_i = \sum_{l \in s_i} y_{i,l}, \; \forall i \tag{54}$$

$$\sum_{l \in N_{in}} y_{i,l} = \sum_{l \in N_{out}} y_{i,l}, \; \forall N, i \tag{55}$$

$$\sum_{l \in L} c_l \sum_{i \in l} y_{i,l} R_i \leq C_L, \; \forall L \tag{56}$$

$$x_i \in \{0, 1\}, \; \forall i \tag{57}$$

$$y_{i,l} \in \{0, 1\}, \; \forall i, l \tag{58}$$

The interpretation of the above problem formulation is as follows. We aim at maximizing the number of routed flows ($\sum_i x_i$) subject to flow i being originated at node s_i (Eqs. (54)) while meeting the flow conservation and capacity constraints (Eqs. (55) and (56)) and imposing that flows cannot be split among different paths (Eqs. (57) and (58)).

The above IP problem can be solved by using standard relaxation techniques [12] which provides as a result an approximation to the optimal single-path routing strategy, which solves the single-path routing problem. The performance of this routing algorithm and the multipath one are evaluated in the next section and compared against other routing algorithms for mesh networks.

5. VALIDATION AND PERFORMANCE EVALUATION

In this section we validate and evaluate the performance of the schemes proposed in this paper. Unless otherwise stated, the experiments with WLAN are based on the 802.11b physical layer.

5.1. Model validation

In order to validate the model presented in Section 3, we performed several experiments and compared the results of our model against those obtained via simulation. The

[11] Note that, while we impose the requirement that the entire flow's throughput is originated at its source, we do not impose any restriction regarding its destination. The reason is that the flow does not have a specific destination and we just want to send it to *any* of the gateways in the mesh. This is a key difference between the classical multi-commodity flow problem and the formulation presented here.

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simulator we used was developed in OMNET++[12] and closely follows the details of the 802.11 protocol.

We simulated a scenario consisting of N flows sharing the same WLAN link, where each flow i sent a throughput proportional to its weight w_i. The weights were allocated according to $\frac{w_{i+1}}{w_i} = \frac{\alpha}{1-\alpha}$ where α is a variable parameter that we used to set different throughput distributions ($\alpha = 0.5$ corresponds to equally distributing throughput among all flows while smaller values of α yield a more uneven distribution).

Table I gives the obtained simulation and analytical results. Results are given for different α values and number of stations for different scenarios. The results show that the analytical values closely follow the ones from simulation, since the error is well below 1% in all cases. We conclude from these results that the analytical model is very accurate.

5.2. Linearized capacity region: Two Flows

In order to evaluate the accuracy of the linearized capacity region mapping to 802.11 proposed in Section 3, we performed the following experiment. We considered two flows sharing the WLAN. In the first scenario, both flows sent at a modulation rate of 11 Mbps (hereafter we refer to this scenario as "homogeneous rates, 11 Mbps"). In the second scenario, one flow is sent at a modulation rate of 11 Mbps and the second one at 1 Mbps ("heterogeneous rates"). The results for the three scenarios are shown in Figs. 4 and 5, respectively. The exact capacity region as given by the exact throughput model presented in Section 3, is plotted with a continuous line. The approximate capacity region as given by the approximation resulting from Eqs. (11)-(13), is plotted with squared dots. Finally the linearized capacity region computed in Section 3.3, with a dotted line. As Figs. 4 and 5 show, the actual capacity region and the approximate one coincide exactly for the homogeneous case, and are quite close for the heterogenous case. This confirms the accuracy of our approximation. Also, the linearized capacity region covers most of the area of the actual capacity region, which means that by linearizing the capacity region we do not waste significant resources. Indeed, the largest deviation between the linearized region and the actual one, which occurs in the axes, does not exceed 10%.

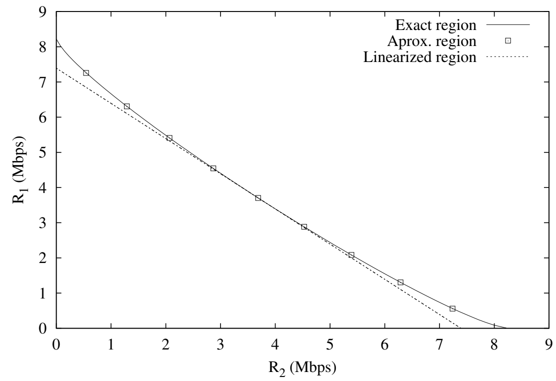

Figure 4. Two Flows, Homogeneous Rates, 11Mbps.

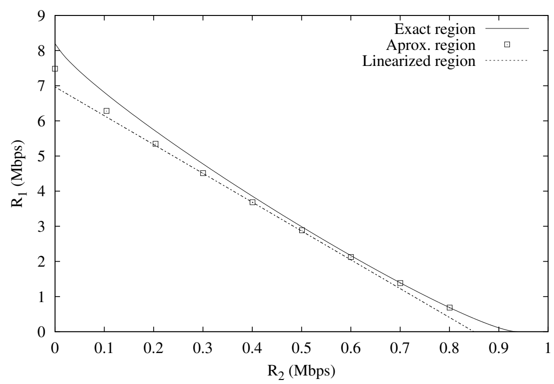

Figure 5. Two stations, Heterogeneous Rates.

5.3. Linearized capacity region: Multiple Flows

In the above experiments only two flows are considered. In order to understand the impact of the number of flows into the behavior of the linearized capacity region, we evaluated the following additional scenarios.

In the first scenario, we considered a varying number of flows sending at the same modulation rate of 11 Mbps ("single rate"). In the second scenario, one fourth of the flows transmit at each of the following modulation rates: 11Mbps, 5.5 Mbps, 2Mbps and 1 Mbps ("multiple rates").

For the each of the above scenarios, we consider a traffic pattern like the one described in Section 5.1, where the traffic distribution depends on a parameter α. For each value of α and number of flows, we depict the total capacity allowed according to the exact capacity region ("exact capacity") and the linearized one ("linearized capacity").

[12] `http://www.omnetpp.org/`

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Scenario	N	Simulation (Mbps)	Analytical(Mbps)
Single Rate (11Mbps), $\alpha = 0.5$	8	7.26	7.22
	16	7.22	7.18
	32	7.20	7.16
Single Rate (11Mbps), $\alpha = 0.1$	8	8.25	8.20
	16	8.25	8.20
	32	8.25	8.20
Two Rates ($\frac{1}{2}$ stations $\{11, 5.5\text{Mbps}\}$), $\alpha = 0.5$	8	5.28	5.26
	16	5.26	5.24
	32	5.25	5.23
Multiple Rates ($\frac{1}{4}$ stations $\{11, 5.5, 2, 1\text{Mbps}\}$), $\alpha = 0.1$	8	7.09	7.10
	16	7.06	7.09
	32	7.02	7.04

Table I. Model validation

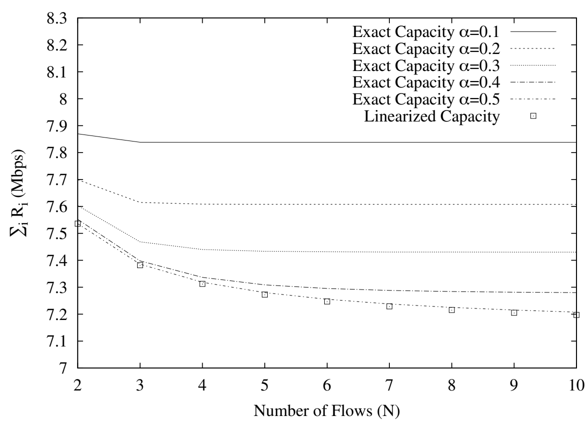

Figure 6. Multiple Flows, Single Rate.

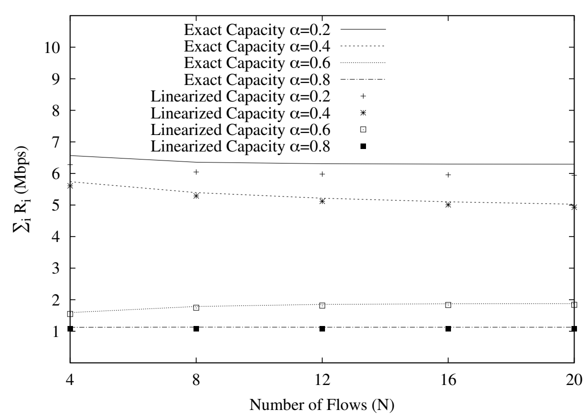

Figure 7. Multiple Flows, Multiple Rates.

The results for the single rate, and multiple rates scenarios are given in Figs. 6 and 7, respectively. Note that because in the single rate scenario the linearized capacity does not depend on α, in Fig. 6 there is only one line drawn for the linearized capacity.The main conclusion that we draw from the above results is that the number of flows has a fairly small impact on the linearized capacity region. Indeed, for the multiple rates scenario the difference between the exact and the linearized capacity does not change noticeably with the number of flows, while for the single rate scenario there is a noticeable change but it is not very significant. These results show that the conclusions given above for two flows also hold for multiple flows.

5.4. Routing: Homogeneous Scenario

The previous experiments have validated the mapping of the linearized capacity region to 802.11. In the following, we evaluate the performance of the proposed routing algorithm based on the linearized capacities. In this section we focus on a homogeneous wireless mesh network consisting of WLAN links only, while in the next section we consider a heterogeneous network consisting of WLAN and WiMAX.

In order to conduct a performance evaluation independent of the chosen topology, we generated multiple random topologies and evaluated the average performance (and its deviation) among all topologies. To generate these random topologies we used the Hyacinth-Laca tool[13], used in several well-known works such as [13] and [14]. We configured this tool to create random topologies with a node count between 40 and 70 nodes (which yields a mean of 55 nodes) spread over an area of 400x400 square meters. Once a topology is available, before performing a routing experiment we need to assign the channels used by each

[13] Available at http://www.ecsl.cs.sunysb.edu/multichannel/

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interface. For this purpose, we used a channel assignment policy that follows a Common Channel Set (CCS) configuration [15–17]. In order to calculate the modulation rate at which each node is able to communicate with its neighbors, we further used the curves of throughput versus distance given in [18].

We evaluated routing performance for 10 gateways and a varying density of source nodes (25, 50 and 100% of the nodes). The metric that we used to evaluate the routing performance is the maximum amount of traffic that can be admitted to the network while providing all flows with the same throughput. For each experiment we generated a set of 35 random topologies, and we provide the average throughput performance and confidence intervals over the throughput resulting from each topology. In order to show the performance improvement resulting from the proposed single-path (SP) and multipath (MP) schemes, we compared them with well-known routing approaches for mesh networks, namely the Expected Transmission Count (ETX) [16], the Expected Transmission Time (ETT) [19] and Shortest Path (ShP).

The results on the routing performance of our approaches (SP and MP) against previous proposals (ETX, ETT and ShP) are given in Fig. 8. From these results, we observe that our two approaches, SP and MP, perform similarly, which means that the performance gain resulting from using multiple paths (MP) is limited. This is because SP can route two flows originating at the same AP through different paths and the additional benefit of splitting individual flows with MP is low. Even more, our two algorithms clearly outperform the other approaches, i.e., they outperform ETX and ETT approximately by a factor of 2 and ShP by a factor of 3.

Among other reasons, the performance improvement of our approach over ETX and ETT are due to the fact that ETT and ETX are additive metrics, and as a result a path consisting of a few rather congested links may be preferred over a longer and less congested path, which harms throughput performance. Furthermore, ETT and ETX do not take into account that the flows that belong to the same wireless link share the same resources; in contrast, our approach considers that allocating throughput to one flow harms the other ones in the same link. We must also consider that by setting the configuration of the MAC parameters according to the algorithm presented in Section 3, our algorithm *jointly optimizes routing decisions and MAC configuration*, which results in an improved performance.

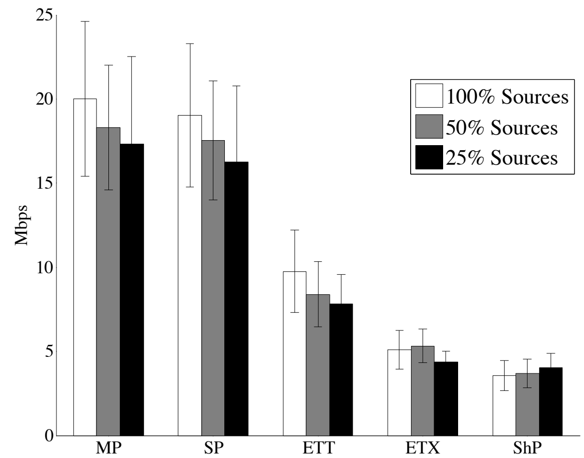

Figure 8. Routing, Homogeneous Scenario, 10 GW nodes.

We conclude from the above results that our method is effective in optimizing throughput performance, making an efficient use of the linearized capacity region and clearly outperforming traditional approaches.

5.5. Routing: Heterogeneous Scenario

In the following we evaluate the performance of our routing algorithm in an heterogeneous WMN based on WLAN and WiMAX technologies. The parameters for WiMAX used in our simulations and the parameters for the linearized capacity region are given in Table II, where c_{uplink} and $c_{downlink}$ are the costs of the uplink and downlink flows, and C is the wireless capacity.

Table II. Linearized Region and physical parameters for WiMAX

Permutation Mode	PUSC	UL	2.2 Mbps
Frame Duration	5 ms	DL	8.4 Mbps
FFT Size	512	C	UL+DL
Modulation	64QAM	c_{uplink}	C/UL
FEC Code	3/4	$c_{downlink}$	C/DL

The scenario we simulated to evaluate the routing performance is the same one as in the previous experiment, with 50% of the nodes acting as sources and 10 gateways. To this scenario, we added a variable number of nodes with WiMAX capabilities (5, 10 and 20 WiMAX nodes, respectively) and an additional node acting as a WiMAX base station.

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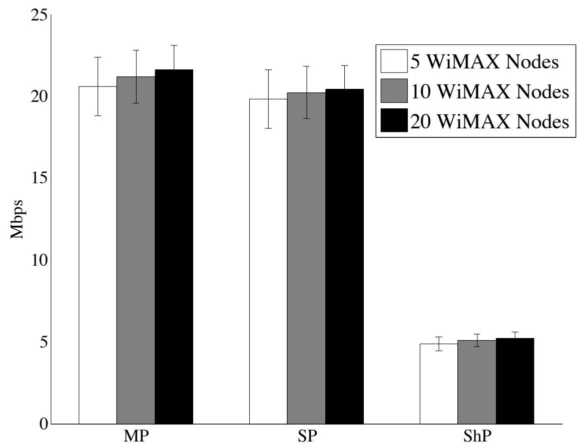

Figure 9. Routing, Heterogeneous Scenario, BS as GW.

Note that the ETT and ETX approaches that we used in the previous section cannot be applied to this scenario as they only work on WLAN-based mesh networks. This way, we only compared the performance of our algorithms against shortest path (ShP), since to our knowledge there are no other approaches in the literature designed for heterogeneous WMNs.

Fig. 9 show the results obtained for the case where the WiMAX base station acts as a gateway. Both our approaches (SP and MP) clearly outperform shortest path (ShP), which confirms that effectiveness of our approach also for the heterogeneous case.

6. RELATED WORK

One of the first and major contributions is the seminal work of Gupta et al. [20] that, based on a geometric analysis, provides an upper bound on the maximum capacity of a network where every node is able to share any portion of the channel it is using with any of its neighbors. Further extensions of this work, e.g., [21, 22] account for the geographic locations and the transmission power of the rate tuples at which a reliable communication is possible. The problem is studied for the case of ad-hoc networks with infrastructure support in [23], and in [24] for the case of WMNs, where different regions of the theoretical maximum capacity are derived depending on the relative ratio of mesh clients, routers and gateways. These works are devoted to the computation of limiting upper bounds on the capacity, and, as such, cannot be used to support optimal routing or perform admission control.

With respect to routing algorithms for WMNs, the first proposals were based on algorithms already available for mobile ad-hoc networks (e.g., hop count); however, given that WMNs significantly differ from MANETS [26] (e.g., because of the infrastructure support), these are far form providing optimal performance. New metrics have been proposed for WMN routing [16, 19, 26, 27], although typically tailored to a particular technology, e.g., IEEE 802.11: ETX [16] is based on the number of attempts to send a frame using lowest-modulation probes; ETT [19] extends it to account for the physical rate and frame length used; ML [27] proposes to find the route with the minimum end-to-end loss probability; while mETX and ENT [26] extend ETX to account not only for average values but also for standard deviations. Our approach significantly differs from these approaches because of two major reasons: i) our proposal is technology agnostic as the routing algorithm is oblivious to the technology supporting a given set of weights and a capacity boundary; and ii) for the particular case of IEEE 802.11, our computation algorithm optimally configures the wireless network, this way minimizing the number of collisions and rendering metrics like ETX that are based on retransmissions useless.

Due to the evolution of wireless networks towards heterogeneity [28], researchers have recently started to address the definition of a QoS-aware metric supporting different technologies [29, 30]. Most of the solutions proposed so far are based on hop counts or nominal bandwidth [31]), this way being unable to support smart routing. A first attempt to use a linear model to optimize the throughput allocation is the work of [32]. However, only the case of 802.11 is considered, and the performance of the MAC protocol is not taken into account –authors assume that the nominal capacity coincides with the achievable capacity of the WLAN.

7. CONCLUSIONS

In this paper we have proposed a novel approach for throughput allocation in heterogeneous mesh networks comprising different wireless technologies.

A key feature of our approach is that it relies in a technology independent interface between the routing layer and the underlying link layers. Such a design decision allows hiding the technology specifics from the underlying links to the routing layer, which is essential for the support of heterogeneous technologies. We note that many previous routing solutions (like e.g. ETT and ETX) rely on technology specific parameters, which makes them unsuitable for heterogeneous mesh networks based on different technologies.

In order to represent the available capacity of a wireless link in a technology independent way, we propose the notion of *linearized capacity region*. The key advantages of our proposal are: i) it relies on a few parameters which yields a simple interface between link layer and routing, ii) by properly adjusting its parameters, the linearized capacity region covers most of the actual capacity region, which provides a high level of efficiency, and iii) by relying on a linear function, it allows for solving network-wide optimization problems at a low computational cost.

The proposed linearized capacity region concept relies on a number of parameters that need to be computed for each of the technologies present in the mesh network. While such a mapping is relatively straightforward for centralized technologies, whose capacity region is already linear, it is more challenging for distributed technologies. In this paper, we have addressed the mapping of the linearized capacity region to a distributed technology like WLAN, by first proposing an algorithm that configures the WLAN to optimize throughput performance, and then linearizing the capacity region resulting from this configuration.

In this paper we have further proposed two routing algorithms on top of the linearized capacity region concept. The first algorithm assumes that flows have the ability of being split over multiple paths (following recent standardization work at the IETF) while the second algorithm is based on traditional single-path flows. The proposed algorithms result from linear and integer programming formulations, respectively, in which the capacity available at each wireless link is given by a linear constraint derived from the linearized capacity region. The approaches proposed in this paper have been extensively evaluated by simulations and show to vastly outperform previous solutions.

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